

String Actuated Curved Folded Surfaces

Martin Kilian

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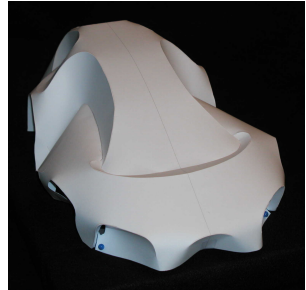
Curved Folding



[Albers, 1927]



[Huffman, 1970s]

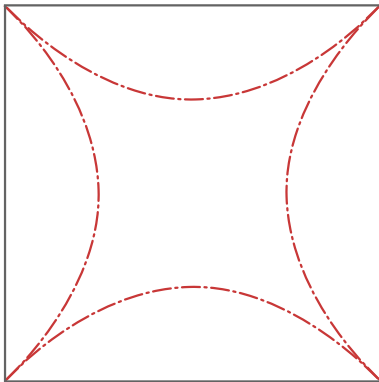


[Epps, 2007]

Models made of paper and PVC.

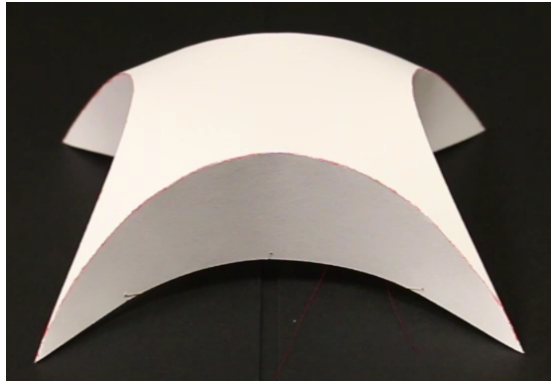
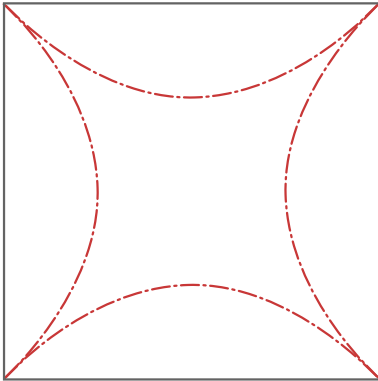
Developable surfaces with curved creases

A simple example



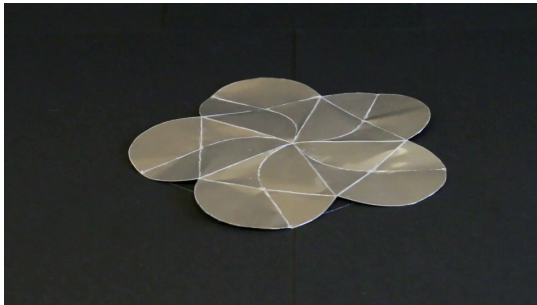
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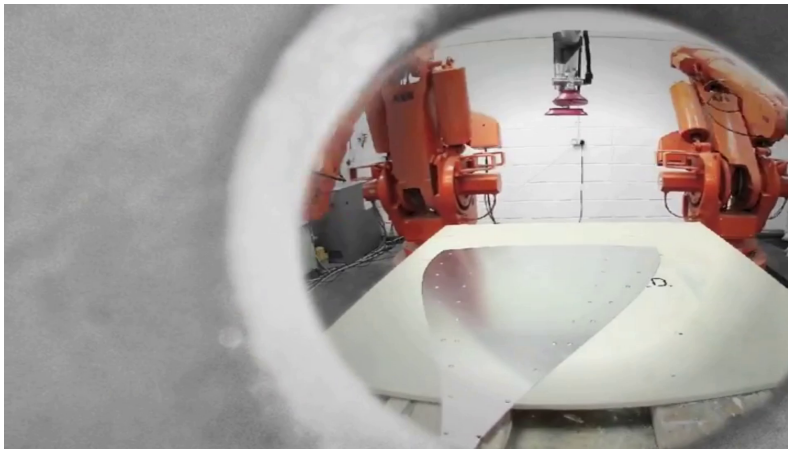


Developable surfaces with curved creases

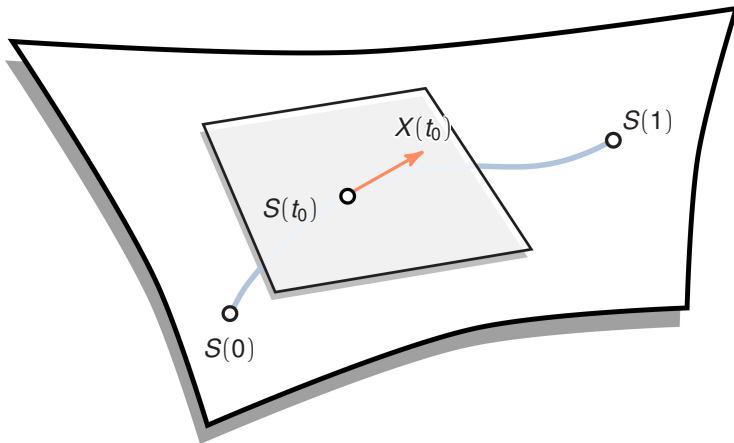
A **static surface** that results from folding a **given** crease pattern.
Folding dynamics have not been considered so far.



Developable surfaces with curved creases

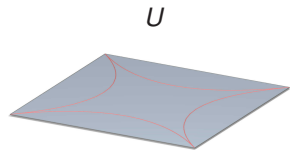


High Level Idea



String actuated folding

Given a **crease pattern** we are looking for a **network of strings** that drives the folding when being pulled. Starting from a **folding sequence**



we compute a set of **actuation points** and a **network of strings** in order to reproduce the given deformation S .

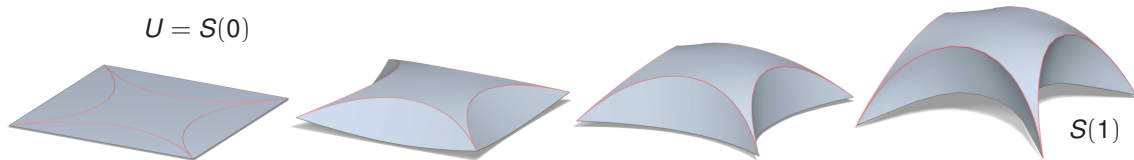
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$$S : [0, 1] \times U \rightarrow \mathbb{R}^3$$

$$(t, \mathbf{u}) \mapsto S(t, \mathbf{u})$$

$U = S(0)$



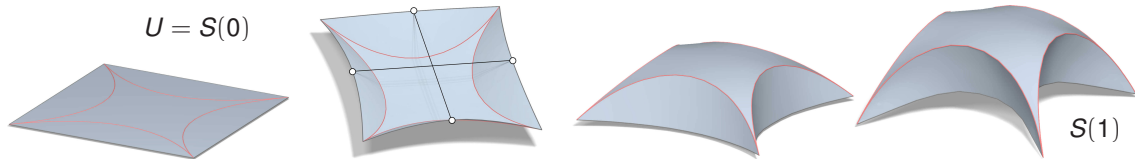
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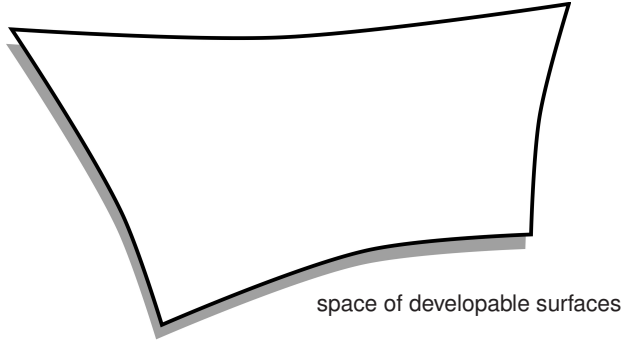
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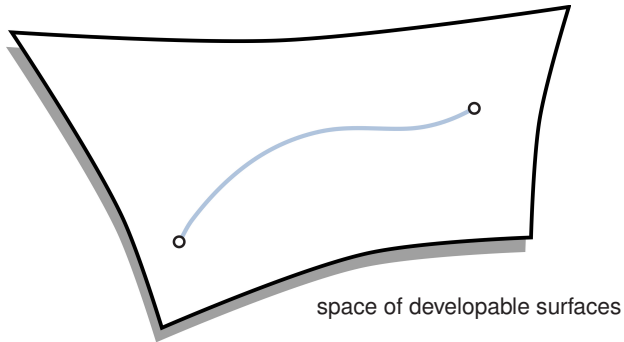


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Analyzing a folding motion

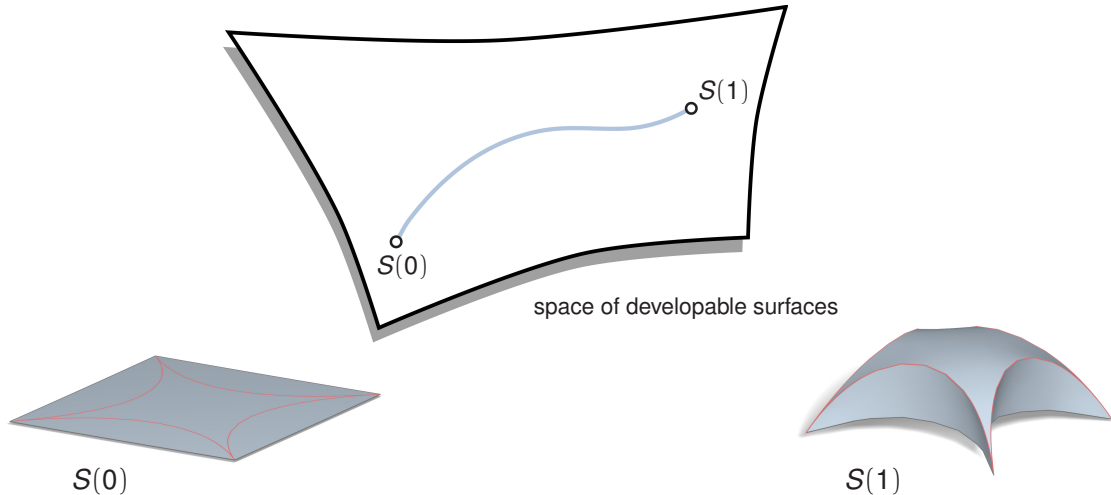


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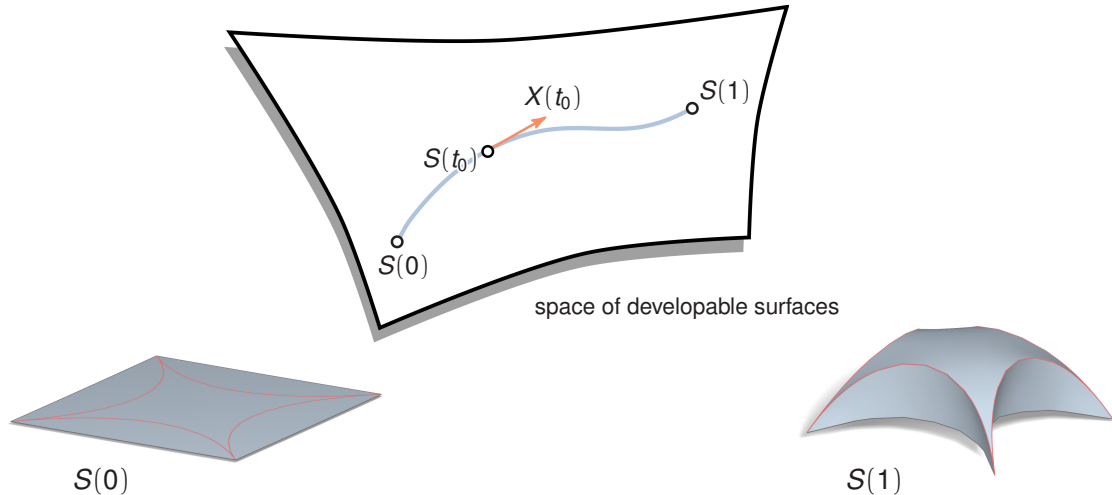


A folding motion S defines a **curve** in the space of developable surfaces.

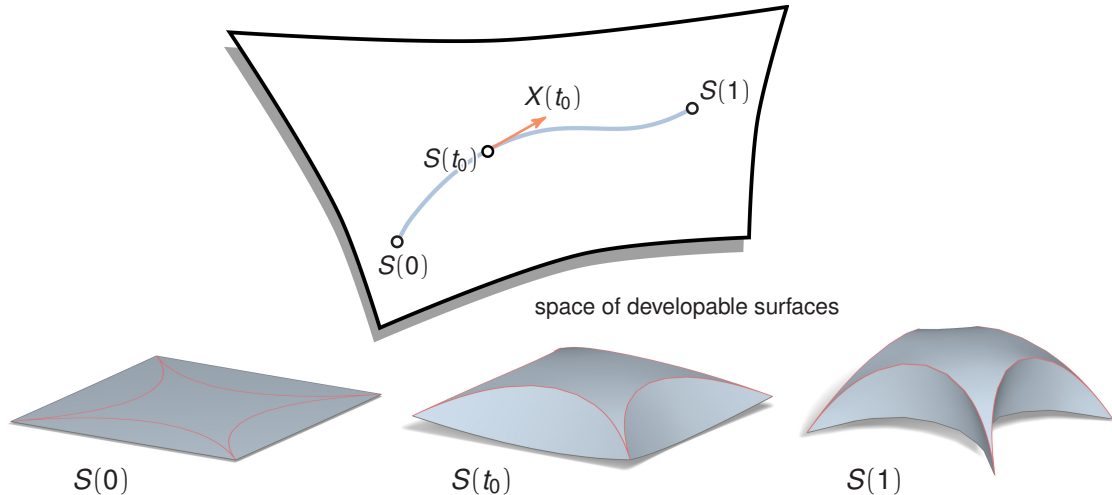
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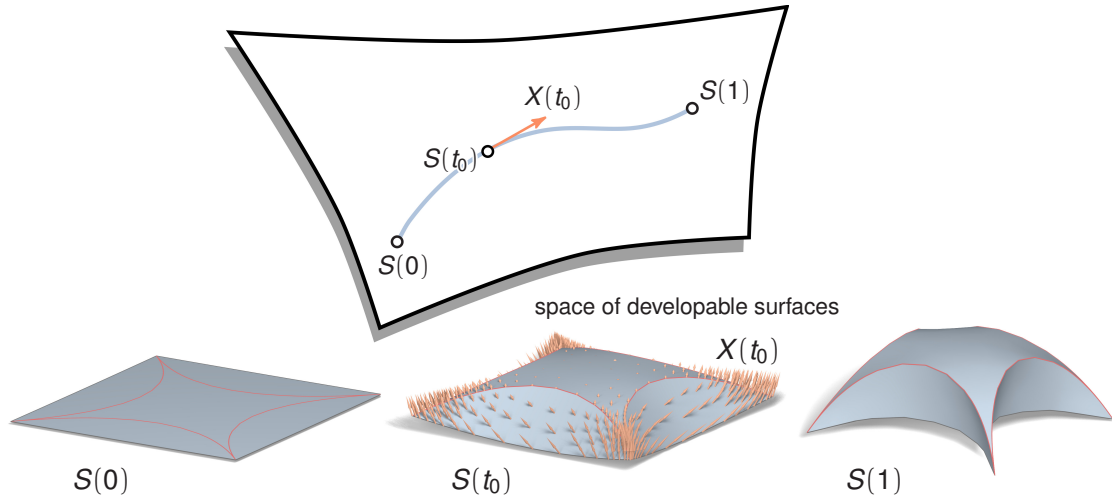
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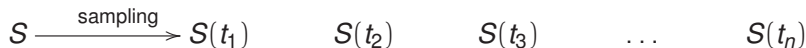


Observation

Any **actuation mechanism** that reproduces the target motion S also has to reproduce the deformation fields $X(t)$, $t \in [0, 1]$.

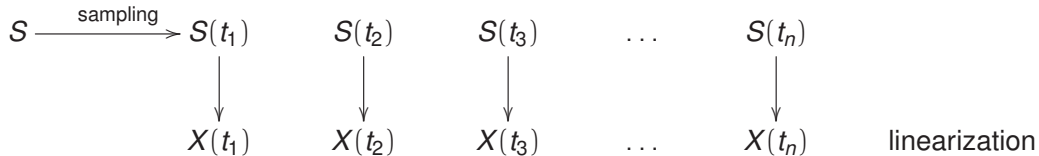
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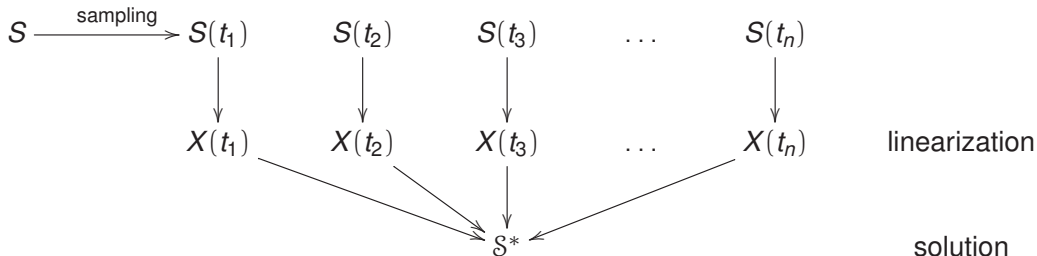
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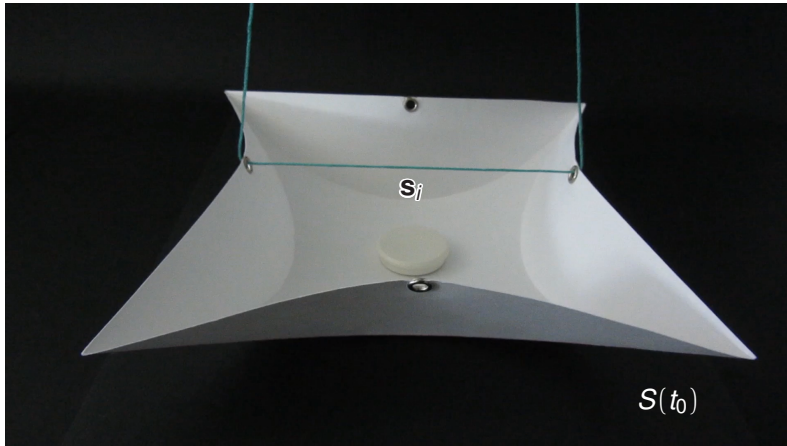
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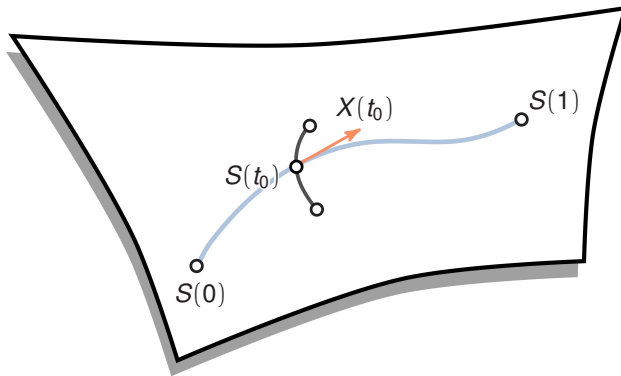


Analyzing a folding motion

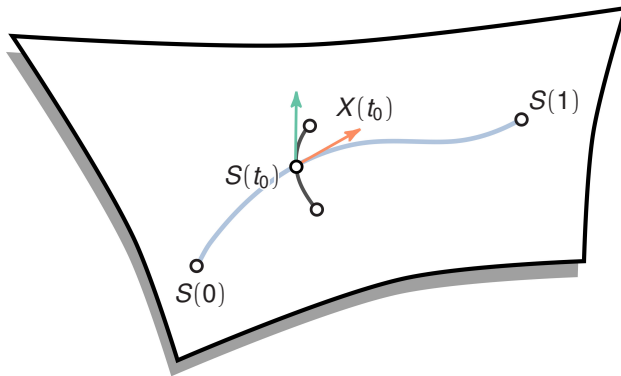
Attach an **arbitrary string** s_i to the **rest shape** $S(t_0)$:



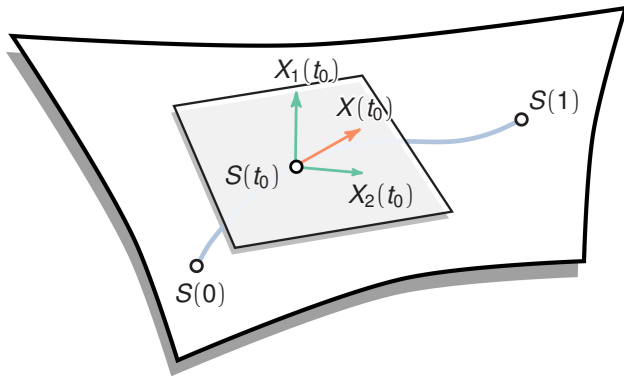
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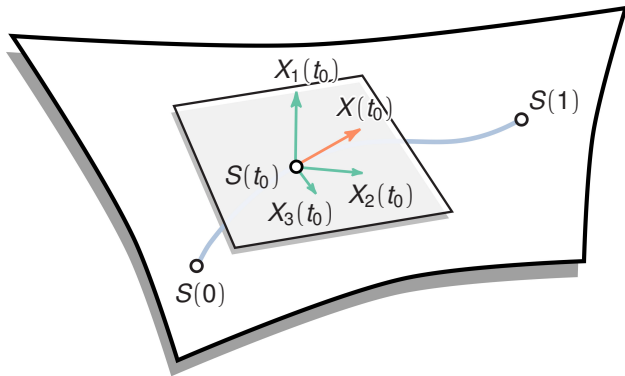


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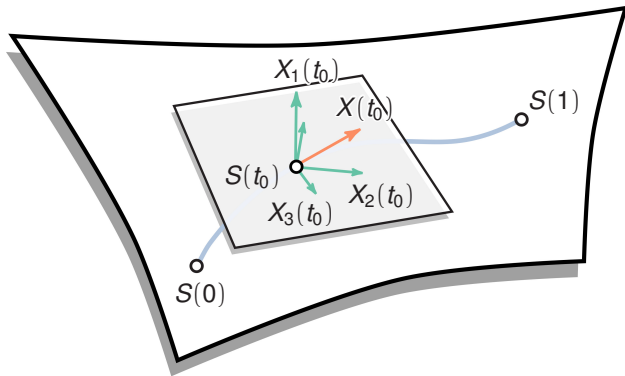
Each $X_i(t_0)$ is the **deformation field** induced by string \mathbf{s}_i , connecting exactly two surface points.

Analyzing a folding motion



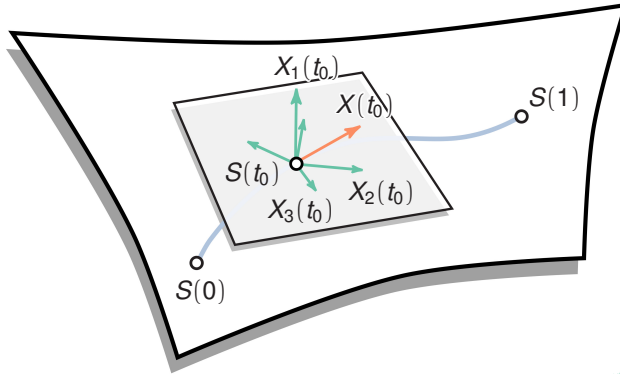
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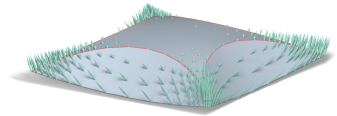


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Analyzing a folding motion



We call the $X_i(t_0)$ **actuation modes**.



Local Problem

At time t_0 find strings \mathbf{s}_i that reproduce $S|_{[t_0-\varepsilon, t_0+\varepsilon]}$.

Local action of \mathbf{s}_i is described by $X_i(t_0)$.

Because of linearization:

joint deformation field of $\{\mathbf{s}_i\}$ is **superposition** of the $X_i(t_0)$.

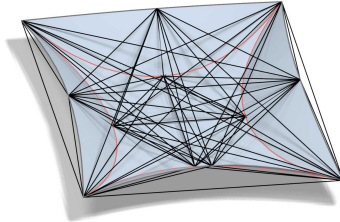
Local Solution

Finding strings \mathbf{s}_i whose joint deformation reproduces $S|_{[t_0-\varepsilon, t_0+\varepsilon]}$ is equivalent to finding actuation modes $X_i(t_0)$ that reproduce $X(t_0)$:

$$\sum \lambda_i X_i(t_0) \approx X(t_0).$$

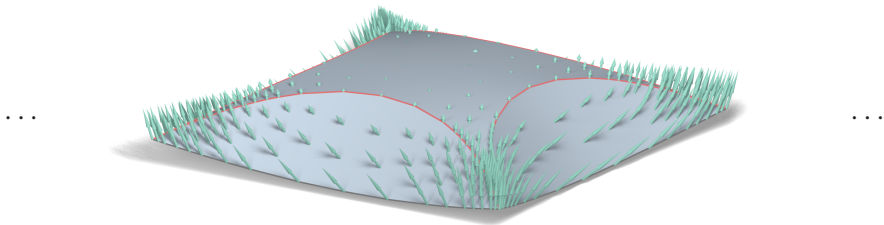
Actuation mode selection

- Sample crease curves to get a set $\mathcal{S} = \{\mathbf{s}_i\}_{i=1}^m$ of **candidate strings**.



Actuation mode selection

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At $S(t_0)$ compute the corresponding modes $X_i(t_0)$.



- Find representation of $X(t_0)$ in terms of the modes $X_i(t_0)$,

$$X(t_0) \approx \sum_{i=1}^m \lambda_i(t_0) X_i(t_0).$$

Actuation mode selection

Compute a **sparse solution** according to

$$\lambda^* = \arg \min_{\lambda} \left[\omega \|\lambda\|_0 + \|X - \sum_{i=1}^m \lambda_i X_i\|^2 \right] \quad (1)$$

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$$\|\lambda\|_0 = \sum_{i=1}^m (\lambda_i \neq 0) \quad ? \quad 1 : 0 \quad (\text{sparsity})$$

$$\|X - \sum_{i=1}^m \lambda_i X_i\|^2 \quad (\text{fitting})$$

Actuation mode selection

Introducing auxiliary variables ξ_i , one can show that

$$\|\lambda\|_0 = \min_{\xi} \sum_{i=1}^m (1 - \xi_i)$$

subject to $\lambda_i \xi_i = 0$ and $0 \leq \xi_i \leq 1$.

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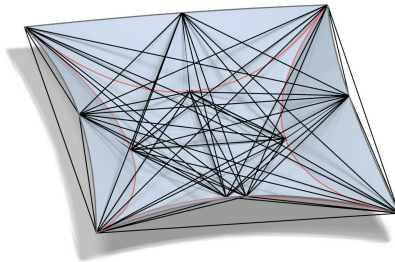
The variables ξ_i act as **indicator** variables:

$$\left[\lambda_i = 0 \iff \xi_i = 1 \right] \quad \text{and} \quad \left[\lambda_i \neq 0 \iff \xi_i = 0 \right]$$

[Feng et al.: **Comp. Formulations of ℓ_0 -norm Optimization Problems**, 2015]

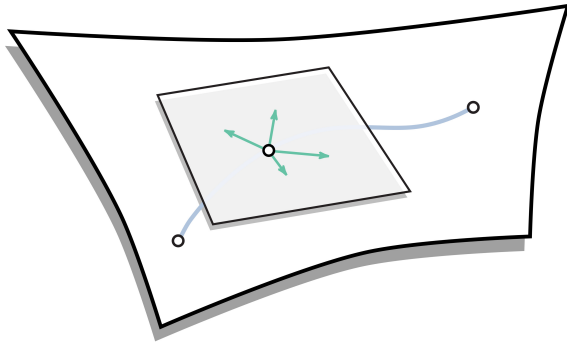
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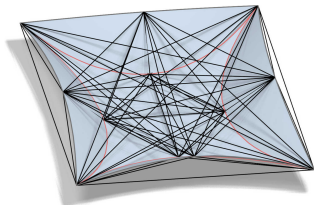
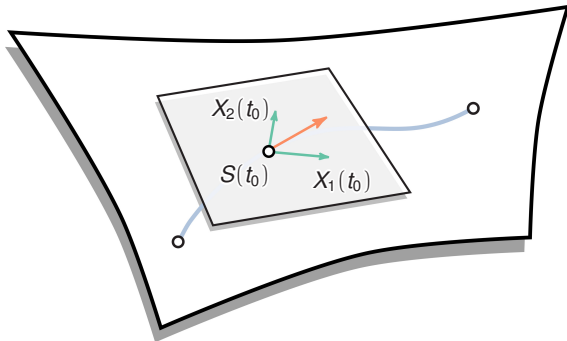
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- 3 solve (1)



[Wächter, Biegler: *On the implementation of a primal-dual interior point filter line search algorithm for large-scale nonlinear programming*, 2006]

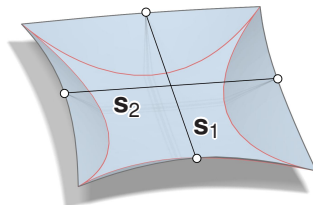
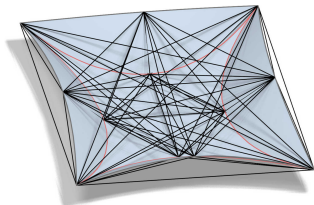
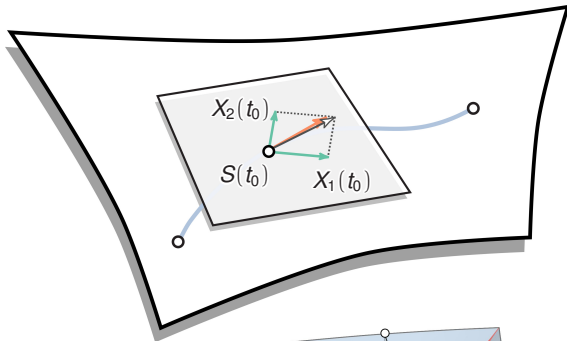
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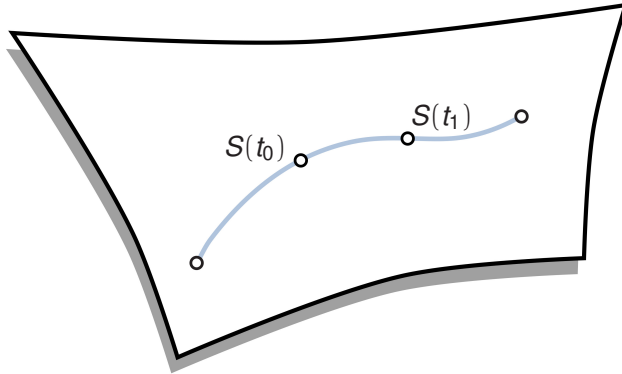


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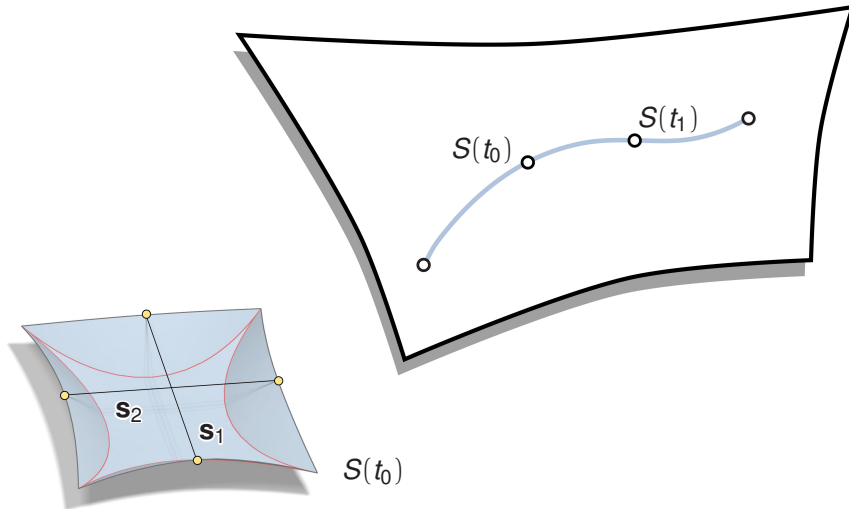
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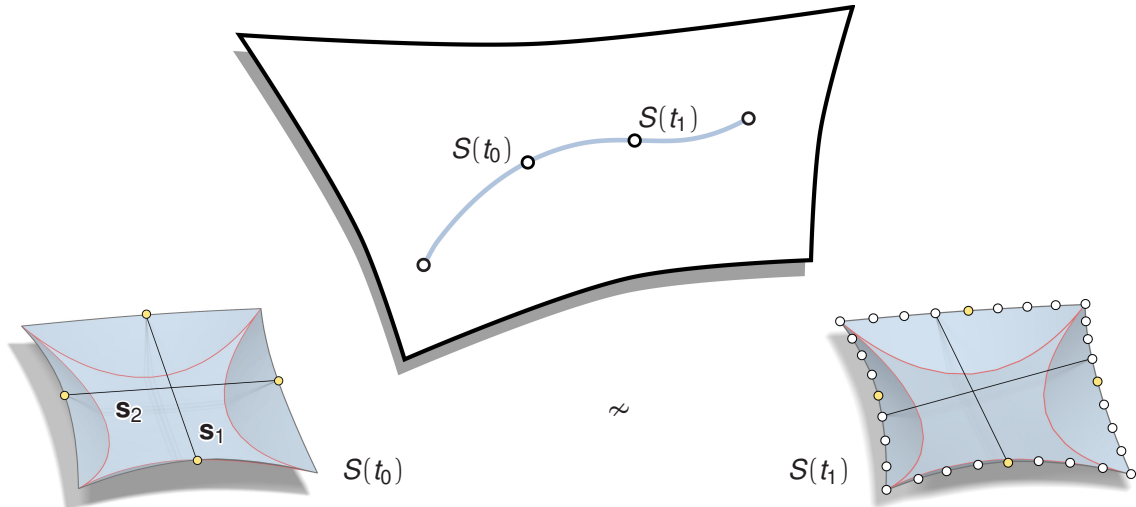
Towards a global solution



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Towards a global solution



Global optimization

Only one ξ_i per \mathbf{s}_i independent of the number of poses $S(t_j)$:

$$\min_{\xi, \lambda} \left[\omega \sum_{i=1}^m (1 - \xi_i) + \sum_{j=1}^n \|X_j - \sum_{i=1}^m \lambda_{ij} X_{ij}\|^2 \right] \quad (2)$$

subject to

$$0 \leq \xi_i \leq 1, \quad \lambda_{ij} \xi_i = 0, \quad i = 1, \dots, m, j = 1, \dots, n.$$

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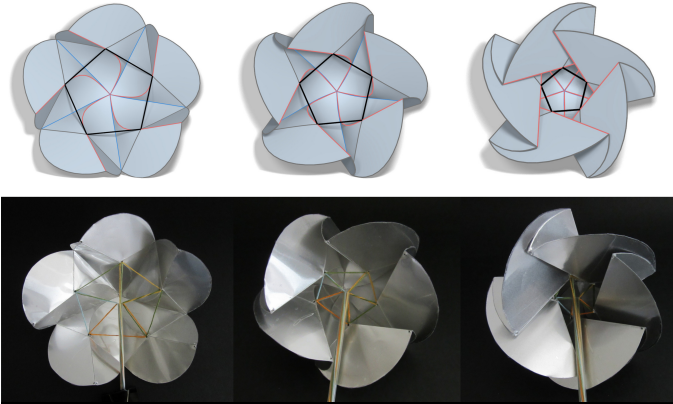
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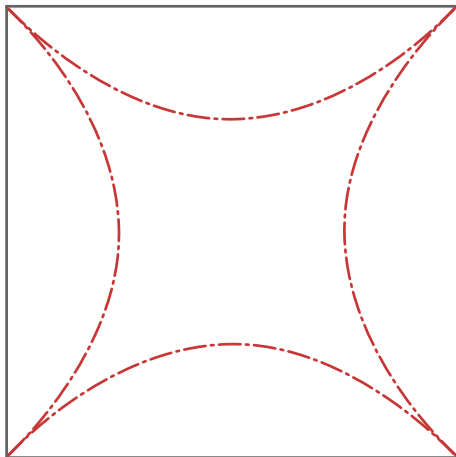
*Have to pay **unit cost 1** to use \mathbf{s}_i .*

Re-using it at other time steps is free!

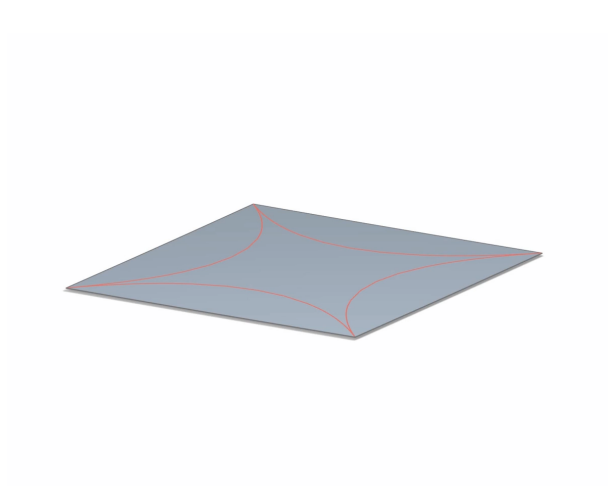
Results



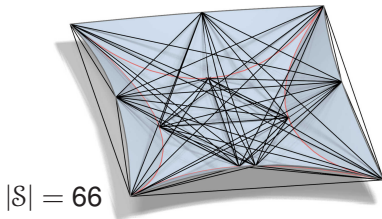
Quad



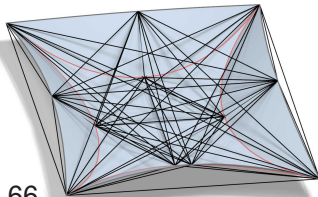
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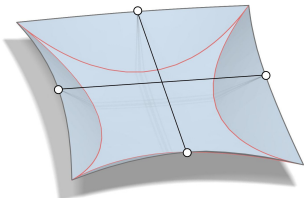
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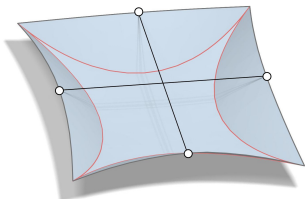
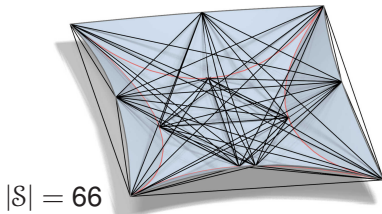


$$|\mathcal{S}| = 66$$

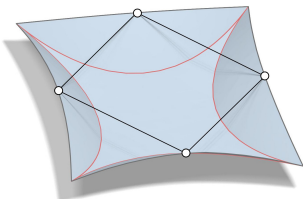


#1

Quad

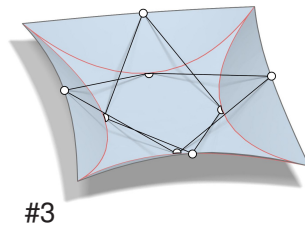
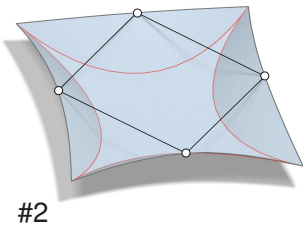
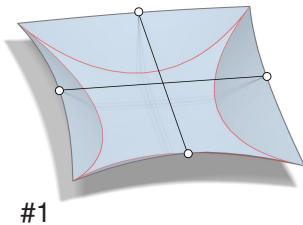
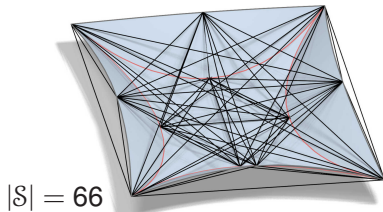


#1

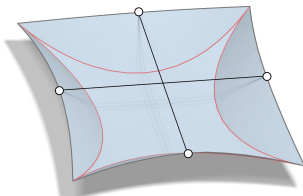
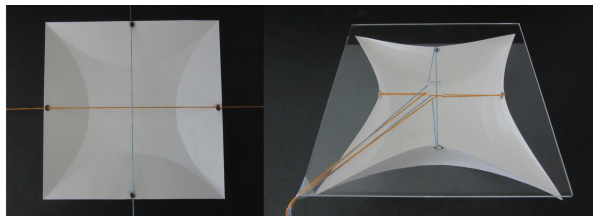
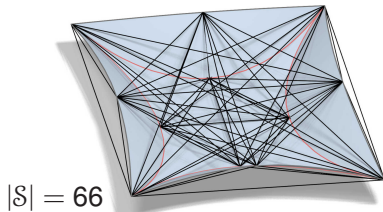


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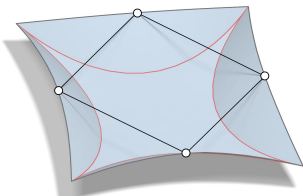
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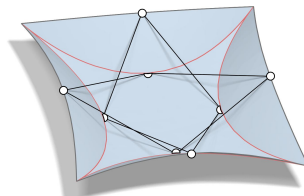
Quad



#1

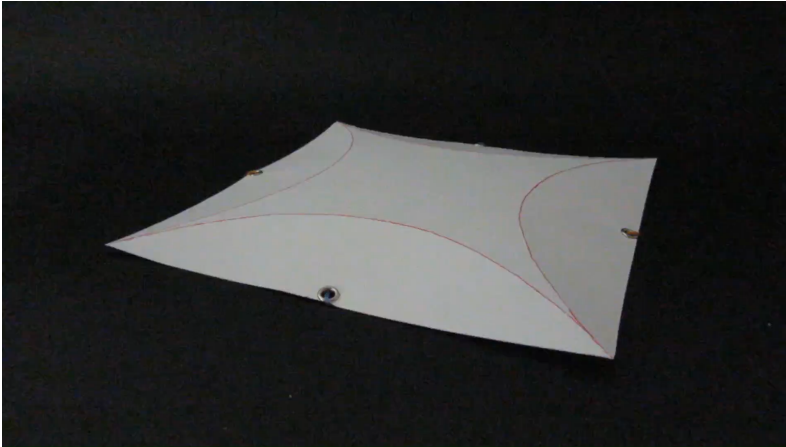


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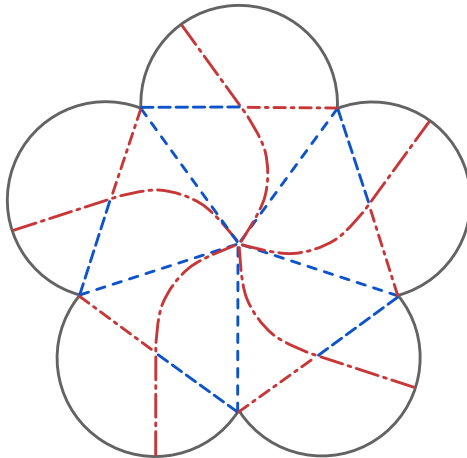


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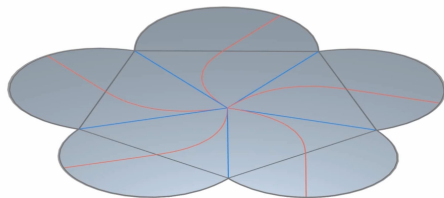
Quad



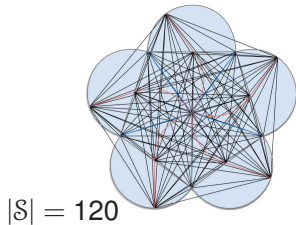
Apricot (Design by J. Mitani)



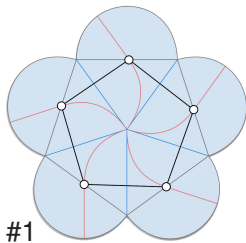
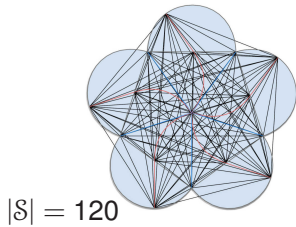
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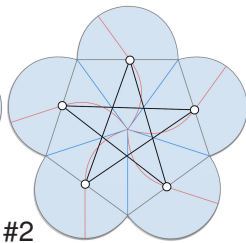
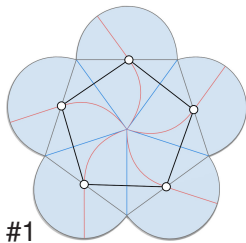
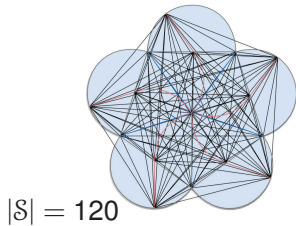
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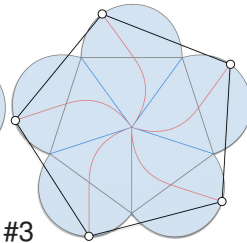
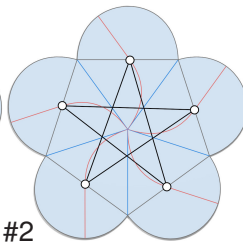
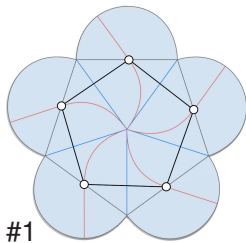
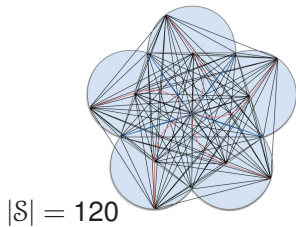
Apricot



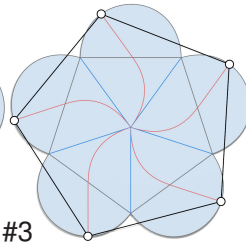
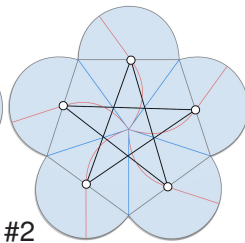
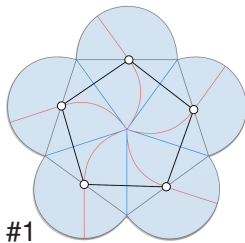
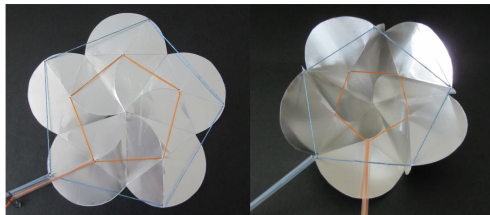
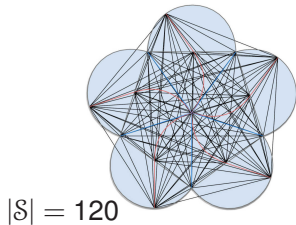
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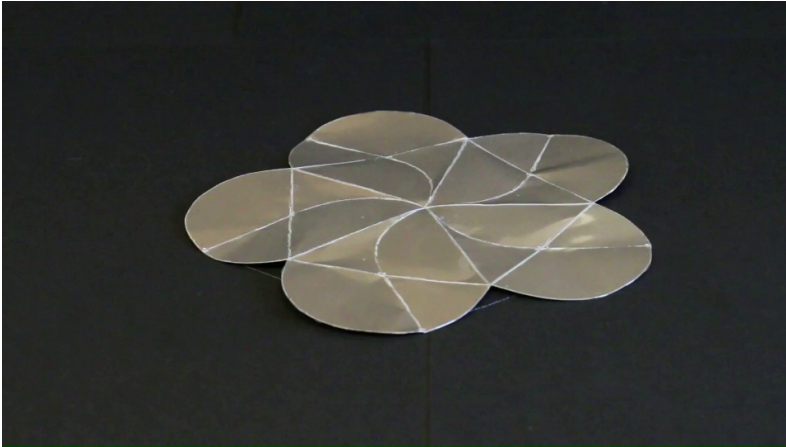
Apricot



Apricot



Apricot



Conclusion

Limitations

- Not every shape is foldable using strings
- Material weight and gravity not considered
- No exterior actuation points

Conclusion

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Future work

- Motion design
- Staged folding (keyframes)

Conclusion

Limitations

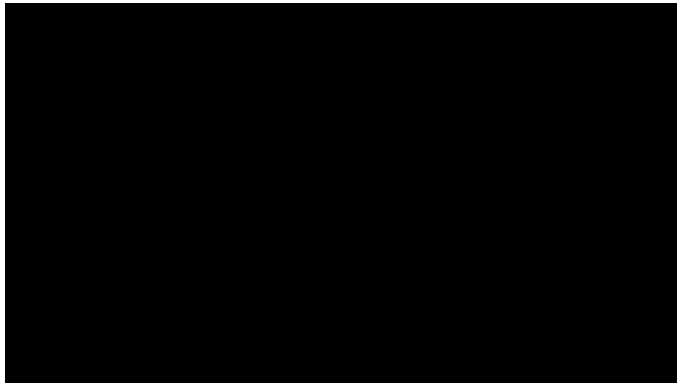
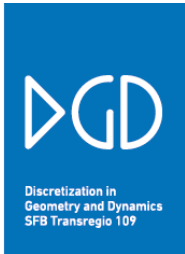
- Not every shape is foldable using strings
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Future work

- Motion design
 - Staged folding (keyframes)
- } Combined pattern, shape, and motion optimization

Thank you!

FWF



<http://geometry.cs.ucl.ac.uk/projects/2017/string-actuated/>